

PARALLEL ALGORITHMS FOR FINANCIAL DERIVATIVES
EVALUATION IN GENERALIZED HESTON MODEL

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Abstract. This paper shows how can be estimated the value of an option if we assume the Heston model on a message-based architecture. We use two methods: first, a Monte Carlo method, then a parallelization of a recurrence obtained from a generalized Merton-Garman equation.

1. INTRODUCTION

From physical models, the following situation has reached acceptance: a financial asset *interest rate* follows a normal law, where the mean is the *drift rate* and the deviation is the *volatility*. This leads to a model that is currently accepted in finance: the model of *geometric Brownian motion*. This model (known as *Black–Scholes–Merton model* in finance and financial engineering) is a stochastic differential equation:

$$dS(t) = mS(t)dt + sS(t)dB(t),$$

where:

- a) $S(t), t \geq 0$ is a stochastic process for the *value of stock*;
- b) m is a static parameter for the *drift rate of return*;
- c) s^2 is a static parameter for the *volatility of stock* ($s \geq 0$);
- d) $B(t), t \geq 0$ is a *standard Wiener process*.

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Another model is assumed by Heston (see [1]) and it consists from two stochastic differential equations: The Heston model corrects some inconsistency of the Black–Scholes–Merton model, for example:

a) in reality, volatility is not a static parameter; it can be used as static value only on short periods (this value will obtain on calibration process, usual with a *statistical estimator*);

b) on long periods, it is possible that interest rate series did not verify a normal law.

A generalization of Heston model is described by the following coupled stochastic differential equations:

$$\begin{aligned}dS(t) &= A(S(t), v(t), t)dt + B(S(t), v(t), t)dB_1(t) \\ dv(t) &= C(S(t), v(t), t)dt + D(S(t), v(t), t)dB_2(t)\end{aligned}$$

where:

- a) $S(t), t \geq 0$ is a stochastic process for value of stock;
- b) $v(t), t \geq 0$ is a stochastic process for volatility of value of stock;
- c) $A(S, v, t), B(S, v, t), C(S, v, t), D(S, v, t)$ are three parametric algebraic functions;
- d) $B_1(t), B_2(t), t \geq 0$ are two *r correlated* standard Wiener processes, i.e.:

$$dB_1(t)dB_2(t) = rdt$$

For Wiener processes, more details can be found in [2].

For the basic Heston model we have:

- a) $A = S(t)m$
- b) $B = S(t)\sqrt{v(t)}$
- c) $C = K(\theta - v(t))$
- d) $D = \xi\sqrt{v(t)}$

where:

- a) m is a drift of rate;
- b) θ is *long run average price volatility*; as t tends to infinity, the expected value of $v(t)$ tends to θ ;
- c) K is the rate at which $v(t)$ reverts to θ ;
- d) ξ is the *volatility of the volatility*; as the name suggests, this determines the variance of $v(t)$.

Note that for $C = D = 0$ we obtain a *static volatility* model (Black–Scholes–Merton):

$$dv(t) = 0.$$

Any financial derivative based on support with price $S(t)$ at time t , with quotation at time t and a value S of support as $V(t, S)$, where $V : [0, T] \times R_+ \rightarrow R_+$ and at maturity time T will generate an interest rate: $\text{payoff} : R_+ \rightarrow R_+$

For example, European options CALL and PUT has payoff functions:

$$\text{payoff}(x) = \max\{0, x - E\}$$

$$\text{payoff}(x) = \max\{0, E - x\}$$

2. PARALLELIZATION OF A MONTE CARLO METHOD FOR OPTION'S EVALUATION

First, we discretize continuous dimension of time. Let us denote:

$$t[k] = t[0] + k\Delta, \quad 0 \leq k \leq N, \quad \Delta = \frac{T - t[0]}{N}$$

where:

a) T is the maturity time of option;

b) N is a number of time units (like days); note that sometimes is used transaction days – in this case, discretization hasn't a constant step.

Because for a standard Wiener process $B(t)$, $t \geq 0$ we can obtain a standard normal random variable series $X[B(t)]$, $t \geq 0$ with:

$$dB(t) = X\sqrt{dt}$$

we can build a simulation step as:

$$A = A(S[k], v[k], t[k])$$

$$B = B(S[k], v[k], t[k])$$

$$C = C(S[k], v[k], t[k])$$

$$D = D(S[k], v[k], t[k])$$

$$S[k+1] = S[k] + A \Delta + B X \text{SQRT}(\Delta)$$

$$v[k+1] = v[k] + C \Delta + D Y \text{SQRT}(\Delta)$$

where X and Y are r correlated. A simple method to generate two r correlated normal values is:

$$X = \text{NORMRAND}()$$

$$Z = \text{NORMRAND}()$$

$$Y = r X + \text{SQRT}(1-r^2) Z$$

A complete simulation for interval $[t_0, T]$ in N step with evaluation of payoff is:

```

FUNCTION simulation()
S = S0
v = v0
t = t0
Δ = (T - t0) / N
FOR k = 1, N
    t = t + Δ
    X = NORMRAND()
    Z = NORMRAND()
    Y = r X + SQRT(1-r2) Z
    A = A(S, v, t)
    B = B(S, v, t)
    C = C(S, v, t)
    D = D(S, v, t)
    SS = S + A Δ + B X SQRT(Δ)
    vv = v + C Δ + D Y SQRT(Δ)
    S = SS
    v = vv
END FOR
RETURN S
END FUNCTION

```

Because for a level of acceptance α , where $0 < \alpha < 1$, a trust interval for $E[S(T)]$ is $[s - a, s + a]$, where

$$s = \sum_{1 \leq k \leq M} simulation()$$

$$a = \frac{F\left(\frac{\alpha}{2}\right)\sigma}{\sqrt{M}},$$

and F is the inverse function for CDF (cumulative distribution function) of standard normal distribution; it means that:

$$\text{Prob}(s - a < E[S(T)] < s + a) = 1 - \alpha$$

or:

$$\text{Prob}(E[S(T)] = s + O(\sqrt{M})) = 1 - a.$$

where *big-O* notation is a *Buchmann-Landau symbol* (see [3]).

Suppose that we have a parallel architecture with message passing protocol (like MPI; for MPI see [5]) that consists in P processors. We broke a simulation job for $P-1$ slaves which will transmit result of $\left\lceil \frac{M}{P-1} \right\rceil$ simulations to master processor:

```

PROGRAM complete_simulation
GLOBAL t0, T, S0, v0, P, N, EPSILON, M
READ t0, T, S0, v0, N, EPSILON, M
P = ProcessorsCount()
IF P ≥ 3 THEN
    CALL parallel_simulation()
ELSE
    CALL serial_simulation()
END IF
END PROGRAM

PROCEDURE parallel_simulation()
LOCAL x, y, Q
IF ProcessorID() = 0 THEN
    FOR i = 1, P-1
        RECV ANY, y
        x = x + y
    ENDFOR
    x = payoff(x / (P-1))
ELSE
    Q = CEIL(M / (P-1))
    x = 0
    FOR i = 1, Q
        x = x + simulation()
    ENDFOR
    SEND 0, x / Q
ENDIF
END PROCEDURE

```

```

PROCEDURE serial_simulation()
LOCAL x
x = 0
FOR i=1, M
    x = x + simulation()
ENDFOR
WRITE payoff(x / M)
END PROCEDURE

```

3. PARALLELIZATION OF NUMERICAL METHOD APPLIED TO GENERALIZED MERTON–GARMAN EQUATION

We recall the following extended Itô's Lemma. For Itô's lemma see [4].

Theorem. Let $S(t)$, $t \geq 0$ and $v(t)$, $t \geq 0$ be two stochastic processes that verify the next differential stochastic equations:

$$dS(t) = A(S(t), v(t), t)dt + B(S(t), v(t), t)dB_1(t)$$

$$dv(t) = C(S(t), v(t), t)dt + D(S(t), v(t), t)dB_2(t)$$

where $B_1(t)$, $t = 0$ and $B_2(t)$, $t = 0$ are two Wiener correlated processes with ρ correlation. If $f(S, v, t)$ is a continuous differentiable function, then:

$$df = \left[f_t + f_S A + f_v C + \frac{f_{SS} B_1^2}{2} + \frac{f_{vv} D^2}{2} + f_{Sv} \rho B_1 D \right] dt + f_S B_1 dB_1 + f_v D dB_2$$

Proof: see [6]

Because we have two sources of risk, we must build a portfolio based on two types of derivatives and support:

- a) one option at value $V(t, S, \sigma)$;
- b) X shares at price $S(t)$;
- c) Y options at value $V_1(t, S, \sigma)$.

Value of this portfolio is:

$$\Pi(t) = V(t, S, \sigma) + X S(t) + Y V_1(t, S, \sigma)$$

After applying extended Itô's lemma for:

$$f_1(S, \sigma, t) = V_1(t, S, \sigma)$$

$$f_2(S, \sigma, t) = V(t, S, \sigma)$$

we can compute:

$$d\Pi(t) = dV(t, S, \sigma) + X dS(t) + Y dV_1(t, S, \sigma)$$

as:

$$d\Pi(t) = \alpha_1 dt + \alpha_2 dB_1(t) + \alpha_3 dB_2(t)$$

Because is a risk-free portofolio, we must have:

$$\alpha_2 = \alpha_3 = 0.$$

But for a risk-free portofolio we have:

$$d\Pi(t) = r\Pi(t) dt.$$

From last two we obtain that (see [6]):

$$\frac{rV + V_s rS + V_t + AV_s + CV_v + \frac{B^2 V_{ss}}{2} + \frac{D^2 V_{vv}}{2} + \rho BDVS_v - AV_s}{DV_v},$$

is an invariant denoted as β , and named as *the market price of volatility risk* (see [7], p8). Rewrite last as a partial differential equation:

$$V_t + rSV_s + (C - D\beta)V_v + \frac{B^2 V_{ss}}{2} + \frac{D^2 V_{vv}}{2} + \rho DVS_v - rV = 0$$

obtain an equation like Black-Scholes equation (more on Black-Scholes equation in [8]). For Heston model, this equation is named as *Merton-Garman equation* (see [9]; [10] p. 41) or *Garman equation*. We name as *generalized Merton-Garman equation* or *generalized Garman equation*.

On boundary we have:

a) value of a null valued stock is null:

$$V(S, v, T) = 0, \forall v, \forall t, 0 \leq t \leq T,$$

b) value at maturity will be payoff():

$$V(S, v, T) = \text{payoff}(S), \forall v, \forall S, S \geq 0.$$

c) for null volatility we have a deterministic solution S^* of differential equation:

$$dS(t) = A(S(t), 0, t) dt$$

and value of option is:

$$V(S, 0, t) = S^*(t), \forall S, \forall t, S \geq 0, 0 \leq t \leq T.$$

In particular case of Heston model:

$$A = r S,$$

solution S^* is:

$$S^*(t) = S(t_0) \exp(r(t - t_0))$$

After discretization on all three axis with ΔS , Δv , Δt denote with:

$$V[i, j, k] = V(j(\Delta S), k(\Delta v), i(\Delta t))$$

generalized Merton–Garman equation can be rewritten as:

$$\begin{aligned} & \frac{V[i, j, k] - V[i-1, j, k]}{\Delta t} + \frac{rj(\Delta S)(V[i, j+1, k] - V[i, j-1, k])}{2\Delta S} + \\ & + \frac{(C - D\beta)(V[i, j, k+1] - V[i, j, k-1])}{2\Delta v} + \frac{B^2(V[i, j+1, k] - 2V[i, j, k] + V[i, j-1, k])}{2(\Delta S)^2} + \\ & + \frac{D^2(V[i, j, k+1] - 2V[i, j, k] + V[i, j, k-1])}{2(\Delta v)^2} + \\ & + \frac{\rho BD(V[i, j+1, k+1] - V[i, j+1, k-1] - V[i, j-1, k+1] + V[i, j-1, k-1])}{(\Delta S)(\Delta v)} - \\ & - rV[i, j, k] = 0 \end{aligned}$$

or in linear form explicit form:

$$\begin{aligned} V[i-1, j, k] = & p V[i, j, k] + q V[i, j+1, k] + r V[i, j-1, k] + \\ & + s V[i, j, k+1] + t V[i, j, k-1] + u V[i, j+1, k+1] - \\ & - u V[i, j+1, k-1] - u V[i, j-1, k+1] + u V[i, j-1, k-1] \\ & - \end{aligned}$$

where:

$$\begin{aligned} p &= (a - 2d - 2e - g) / a \\ q &= (d + b) / a \\ r &= (d - b) / a \\ s &= (e + c) / a \\ t &= (e - c) / a \end{aligned}$$

$$\begin{aligned}
u &= f / a \\
a &= (\Delta S)^2 (\Delta v)^2 \\
b &= \frac{1}{2} r j (\Delta t) (\Delta S)^2 (\Delta v)^2 \\
c &= \frac{1}{2} (C - D \beta) (\Delta t) (\Delta S)^2 (\Delta v) \\
d &= \frac{1}{2} B^2 (\Delta t) (\Delta v)^2 \\
e &= \frac{1}{2} D^2 (\Delta t) (\Delta S)^2 \\
f &= \rho B D (\Delta t) (\Delta v) (\Delta S) \\
g &= r (\Delta t) (\Delta v)^2 (\Delta S)^2
\end{aligned}$$

A parallel solution for a PRAM (see [11]) architecture was presented in [12]. We will build a parallel solution for a message passed architecture with a diagonalization method of domain splitting in order to reduce the number of messages (like in [13], where communication was halved). Suppose that our architecture has Q processors, where one will be master and $Q-1$ will be slaves.

Each slave will process a prism obtained from cube

$$\text{CUBE} = [0, M] \times [0, N] \times [0, P],$$

where M is the number of epochs (time), $N-1$ is number of intermediar points on S axis, with $N (\Delta S)$ enough biger to cover values, respectively $P-1$ for v axis, after split it in $Q-1$ prisms:

$$\text{PRISM}[x] = [0, M] \times \{(y, z) : f(x-1) \leq y + z \leq f(x)\}$$

where:

$$f(x) = \text{FLOOR}((N + P + 1) x / (Q - 1)).$$

Because processing of $\text{PRISM}[x]$ need first 2 left layer from $\text{PRISM}[x+1]$ and last 2 right layer from $\text{PRISM}[x-1]$, processors must interchange some values $M-1$ times. Parallel algorithm is:

PROGRAM simulation

$Q = \text{ProcessorsCount}()$

$x = \text{ProcessorID}()$

IF $x = 0$ THEN

$Z = 0$

WHILE $Z < Q - 1$ DO

RECV x, I, J, K, a

IF $I = J = K = 0$ THEN

$Z = Z + 1$

```

ELSE
  V[I, J, K] = a
ENDIF
END WHILE
ELSE
  FOR j = 0, N
    FOR k = 0, P
      V[M, j, k] = payoff(j (ΔS))
    ENDFOR
  ENDFOR
  FOR I = M, 1, -1
    // border
    FOR k = 0, P
      V[i-1, 0, k] = 0
    ENDFOR
    FOR j = 0, N
      V[i-1, j, 0] = S*((i-1) (Δt))
    ENDFOR
    // first 2 layers
    If x <> 1 THEN
      RECV x-1, {V[i,j,k]} | j+k=f(x-1)-1 OR j+k=f(x-1)-2}
    ENDIF
    FOREACH j,k WITH j+k=f(x-1) OR j+k=f(x-1)+1
      CALL compute_one()
    ENDFOR
    If x <> 1 THEN
      SEND x-1, {V[i,j,k]} | j+k=f(x-1) OR j+k=f(x-1)+1}
    ENDIF
    // other internal layers
    FOREACH j,k WITH f(x-1)+2<=j+k<=f(x)-2
      CALL compute_one()
    ENDFOR
    // last 2 layers
    If x <> Q-1 THEN
      RECV x+1, {V[i,j,k]} | j+k=f(x)+1 OR j+k=f(x)+2}
    ENDIF
    FOREACH j,k WITH j+k=f(x) OR j+k=f(x)-1
      CALL compute_one()
    ENDFOR
  
```

```

If x <> Q-1 THEN
  SEND x+1, {V[i,j,k]} | j+k=f(x) OR j+k=f(x)-1}
ENDIF
ENDFOR
SEND 0, 0, 0, 0, 0
ENDIF
END

```

```

PROCEDURE compute_one()
IF j = 0 OR k = 0 THEN
  RETURN
ENDIF
A = A(j (ΔS), k (Δv), i (Δt))
B = B(j (ΔS), k (Δv), i (Δt))
C = C(j (ΔS), k (Δv), i (Δt))
D = D(j (ΔS), k (Δv), i (Δt))
a = (ΔS)2 (Δv)2
b = ½ r j (Δt) (ΔS)2 (Δv)2
c = ½ (C - D β) (Δt) (ΔS)2 (Δv)
d = ½ B2 (Δt) (Δv)2
e = ½ D2 (Δt) (ΔS)2
f = ρ B D (Δt) (Δv) (ΔS)
g = r (Δt) (Δv)2 (ΔS)2
p = (a - 2 d - 2 e - g) / a
q = (d + b) / a
r = (d - b) / a
s = (e + c) / a
t = (e - c) / a
u = f / a
V[i-1,j,k] = p V[i,j,k] + q V[i,j+1,k] + r V[i,j-1,k]
+ s V[i,j,k+1] + t V[i,j,k-1] + u V[i,j+1,k+1]
- u V[i,j+1,k-1] - u V[i,j-1,k+1] + u V[i,j-1,k-1]
SEND 0, i-1, j, k, V[i-1,j,k]
END

```

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